

# Security of Online Social Networks

## Social Network Analysis II

Lehrstuhl IT-Sicherheitsmanagement

Universität Siegen

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# Overview Lesson 07

## Clustering

Clustering Quality

$k$ -means

## Graph Matching

De-Anonymizing Social Networks

# Clustering

# Overview Clustering

- ▶ Classification of Samples
- ▶ Profiling
- ▶ Finding “Groups”
- ▶ “Hidden” similarities
- ▶ e.g. people most likely inclined to buy X
- ▶ e.g. behavioral anomalies

# Clustering Example

- ▶ Clustering: Partition  $\Pi$  of a set
- ▶ Partitioning of a birthday cake (without measure)
- ▶ e.g. determining “supporters” (Example)

# Clustering Quality Measures

## Quality Criteria

- ▶ External (absolute)
  - ▶ using information not available to algorithm
  - ▶ “Cake-Criteria”: Equivalent Partition-Size
- ▶ Internal (relative)
  - ▶ allows for  $Q(\Pi') < Q(\Pi'')$
  - ▶ e.g. silhouette coefficient

## Silhouette Coefficient I

Internal Quality Criterion for Quality of Placement of Sample  $s(a)$

Clustering  $\Pi = \{\pi_1, \dots, \pi_k\}, a \in \pi \subseteq \mathcal{A}$

$$I(a) := \frac{1}{|\pi_i|} \sum_{x \in \pi_i} d(x, a)$$

$$O_j(a) := \frac{1}{|\pi_j|} \sum_{x \in \pi_j} d(x, a)$$

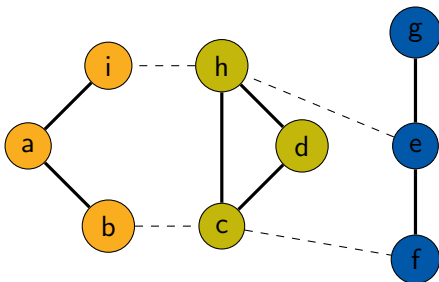
$$O(a) := \min\{O_1(a), \dots, O_{i-1}(a), O_{i+1}(a), \dots, O_k(a)\}$$

$$s(a) := \frac{O(a) - I(a)}{\max\{O(a), I(a)\}}$$

## Silhouette Coefficient II

- ▶ Partition:  $mean\{s(a)\}_{a \in \pi}$
- ▶ Clustering:  $mean\{s(\pi)\}_{\pi \in \Pi}$

Example:





	I	O	s
a	2/2	$\min\{7/3, 10/3\}$	$4/3 * 3/7 = 4/7$
b	3/2	$\min\{5/3, 9/3\}$	$1/6 * 2/3 = 1/9$
i	3/2	$\min\{5/3, 8/3\}$	$1/6 * 2/3 = 1/9$
c	2/2	$\min\{6/3, 6/3\}$	$1 * 2 = 2$
d	2/2	$\min\{7/3, 7/3\}$	$4/3 * 3/7 = 4/7$
h	2/2	$\min\{5/3, 5/3\}$	$2/3 * 3/5 = 2/5$
e	2/2	$\min\{5/3, 8/3\}$	$2/3 * 3/5 = 2/5$
f	3/2	$\min\{5/3, 8/3\}$	$1/6 * 2/3 = 1/9$
g	3/2	$\min\{8/3, 11/3\}$	$7/6 * 3/8 = 7/16$
		$\pi_1$	$\pi_2$
s	$50/189 \approx 0.26455$	$104/105 \approx 0.9905$	$683/2160 \approx 0.3162$
$s(\Pi) \approx 0.523751$			

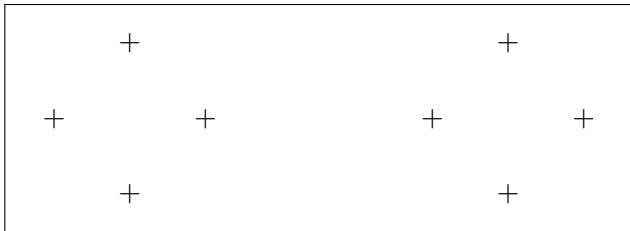
## $k$ -means

Basic  $k$ -means:

- ▶ samples  $\mathcal{A}$
  - ▶  $k$  centroids  $c_1, \dots, c_k \in \mathcal{C}$
1. elements closest to a centroid belong to that partition
    - ▶  $\pi_l = \{a : a \in \mathcal{A}, d(c_l, a) \leq d(c_i, a), \text{ for } l \in \{1, \dots, k\}\}$
  2. adapt centroids to partition
    - ▶  $c(\pi_i) = \arg \min_{x \in \mathcal{C}} \left\{ \sum_{a \in \pi_i} d(x, a) \right\}$
  3. iterate until sufficient quality

## Example $k$ -means

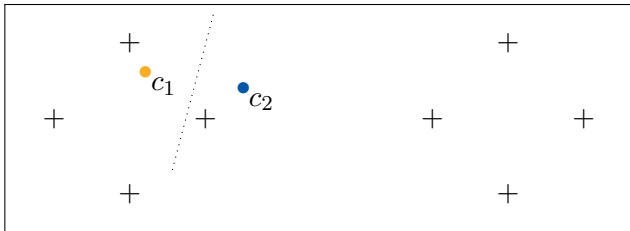
$$k = 2$$



$$\pi_l = \{a : a \in \mathcal{A}, d(c_l, a) \leq d(c_i, a), \text{ for } l \in \{1, \dots, k\}\}$$

## Example $k$ -means

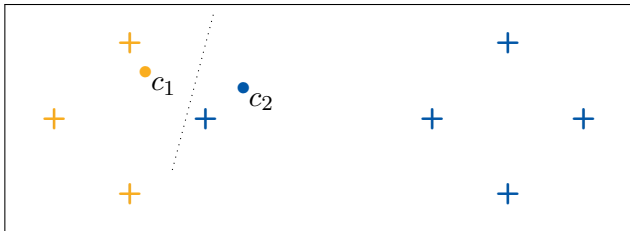
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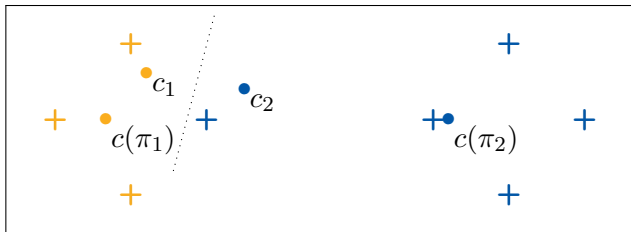
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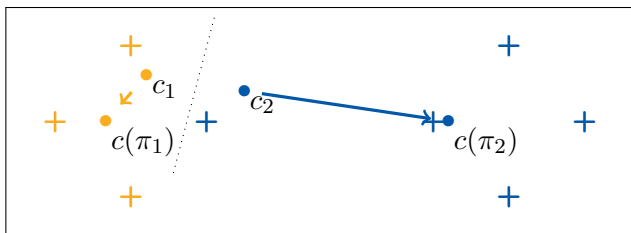
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$$c(\pi_i) = \arg \min \left\{ \sum_{x \in \mathcal{C}} d(x, a) \right\}_{a \in \pi_i}$$

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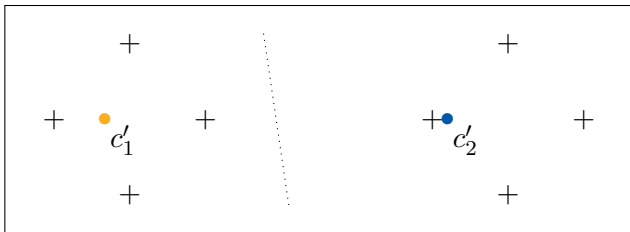
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## Example $k$ -means

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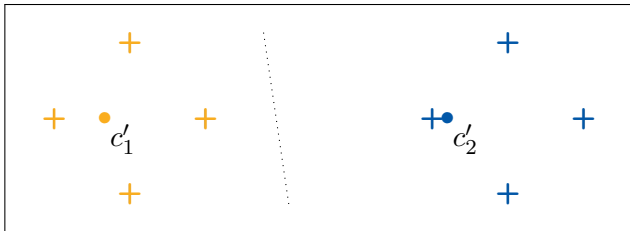


Finish?



## Example $k$ -means

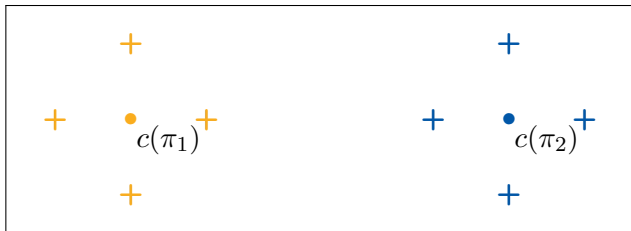
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## Example $k$ -means

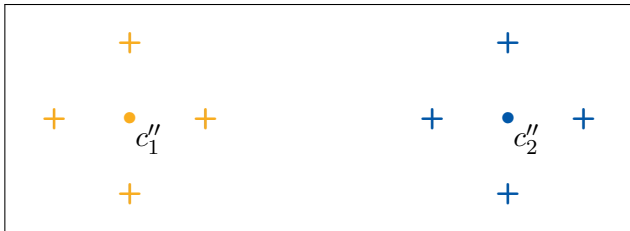
$$k = 2$$



Finish?

## Example $k$ -means

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Finish?

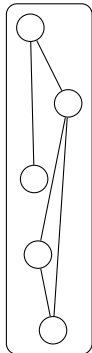
# Clustering Algorithms

- ▶ Halting Problem
- ▶ Sensitive to initial conditions
- ▶ Local Maximum
- ▶ Batch vs. Incremental
- ▶ Large Collection of Algorithms:
  - ▶  $k$ -means (Centroid-based)
  - ▶ Density-based
  - ▶ Distribution-based
  - ▶ Neuronal Networks
  - ▶ Hierarchical Clustering (e.g. greedy least distance)
  - ▶ ...

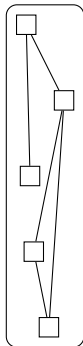
# Graph Matching

## Graph Matching

SocNet1

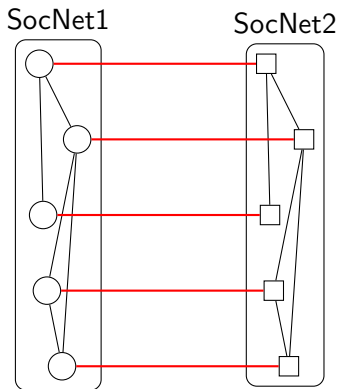


SocNet2



- ▶ Re-identification
- ▶ Imagine this with [twitter](#) and [facebook](#)
- ▶ Hypothesis Space Size?  $n!$
- ▶ Improve your chances with seeds
- ▶ Iterate Nodes ( $n$ )

# Graph Matching



- ▶ Re-identification
- ▶ Imagine this with [twitter](#) and [facebook](#)
- ▶ Hypothesis Space Size?  $n!$
- ▶ Improve your chances with seeds
- ▶ Iterate Nodes ( $n$ )

## De-Anonymizing Social Networks (2009)

- ▶ Narayanan, Shmatikov
- ▶ Sanitised  $\mathcal{S}_{SAN}$ : published, anonymised graph
- ▶ Auxiliary  $\mathcal{S}_{AUX}$ : different attributed graph
- ▶ Objective: Match sanitised onto auxiliary network
- ▶ Procedure:
  1. Seed identification
  2. Propagation



## Model & Definitions

- ▶ Social Network  $\mathcal{S}$ :
  - ▶  $G = (V, E, \mathcal{X}, \mathcal{Y})$ : Nodes, Edges, Attributes
- ▶ Sanitisation/Data Release
  - ▶ Removing Attributes, Nodes, Edges
  - ▶ Fake Edges, ...
  - ▶ Sanitised graph  $\mathcal{S}_{SAN}$
- ▶ Attack Scenarios
  - ▶ Global Surveillance
  - ▶ Abusive Marketing
  - ▶ Phishing and Spamming
  - ▶ Targeted Deanonimisation

# Attacker Model

- ▶ Attacker knows second graph  $\mathcal{S}_{AUX}$
- ▶ Probabilistic Auxiliary Information
  - ▶ Edge attribute “friendship type”
  - ▶ does not cover xor-attributes well

## Seed identification

- ▶ Assumption:
  - ▶ attacker knows of  $k$ -clique in  $\mathcal{S}_{SAN}$ ,  $\mathcal{S}_{AUX}$
- ▶ Input:
  1. target graph ( $\mathcal{S}_{SAN}$ )
  2.  $k$  seed nodes in  $\mathcal{S}_{SAN}$
  3.  $k$  node degrees
  4.  $\binom{k}{2}$  common neighbour counts
- ▶ Output: matching clique in  $\mathcal{S}_{SAN}$
- ▶ error factor  $1 + \epsilon$  on degree and common neighbour
- ▶ return after first match improves efficiency

# Propagation

- ▶ Input:
  - ▶  $G_1 = (V_1, E_1), G_2 = (V_2, E_2)$
  - ▶ partial seed  $\mu_S : V_1 \rightarrow V_2$
- ▶ Output: mapping  $\mu$
- ▶ Iteration:
  1. take arbitrary unmapped  $u \in V_1$
  2. compute for each  $v \in V_2$ : |neigh of  $u$  mapped to neigh of  $v$ |
  3. if  $> thresh$  include  $(u, v)$  in mapping
- ▶ Additional Heuristics:
  - ▶ Minimum eccentricity (of score)
  - ▶ Edge directionality (sum in- and outgoing)
  - ▶ Bias of high-degree nodes (divide score by  $\sqrt{deg(v)}$ )
  - ▶ Revisit nodes (more errors in empty mapping)
  - ▶ Reverse match (compute with reversed inputs)

# Experiments I

- ▶ 100K-nodes graph from real social network
- ▶ random perturbation of auxiliary
- ▶ Number of seeds crucial for propagation
  - ▶ low threshold e.g. 30
  - ▶ overlap node: 25%, edge: 50%  $\Rightarrow \sim 75\%$  re-identified
- ▶ Propagation robust against perturbation

## Experiments II




- ▶ Twitter, Flickr, LiveJournal crawlings of 2007/'08

Network	Nodes	Edges	Av. Deg.
Twitter	224K	8.5M	37.7
Flickr	3.3M	53M	32.2

[Narayanan/Shmatikov 2009]

- ▶ 150 nodes seed
- ▶ name-based ground truth mapping (error < 5%)
- ▶ 30.8% re-identified, 12.1% incorrect, 57% not identified
- ▶ 41% of incorrect are within dist 1 of true mapping
- ▶ 55% of incorrect are within similar geo-location
- ▶ 27% of incorrect are completely erroneous

# Literatur I

-  J. Kogan, *Introduction to Clustering Large and High-Dimensional Data*. Cambridge University Press, 2007.
-  A. K. Jain and R. C. Dubes, *Algorithms for Clustering Data*, ser. Prentice Hall Advanced Reference Series : Computer Science. Prentice Hall, March 1988.
-  A. Narayanan and V. Shmatikov, "De-anonymizing social networks," in *Proceedings of the 30th IEEE Symposium on Security and Privacy (S&P 2009), 17-20 May, Oakland, California, USA*. IEEE Computer Society, 2009, pp. 173–187. [Online]. Available: [http://www.cs.utexas.edu/~shmat/shmat\\_oak09.pdf](http://www.cs.utexas.edu/~shmat/shmat_oak09.pdf)